SNIA DEVELOPER CONFERENCE SNIA DEVELOPER CONFERENCE BY Developers FOR Developers

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Efficient Media Utilization Across Dissimilar Cloud Storage Systems

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- Background Information
- Effects of Layering
- New Media Types
- New Object Store
- New Development Platform
- Conclusion

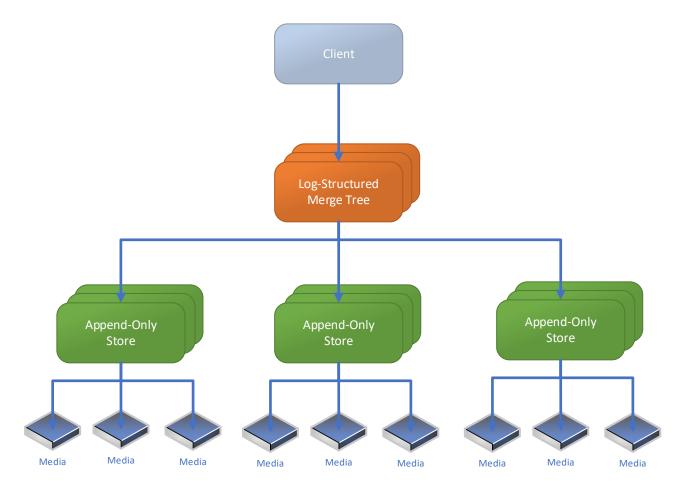




# **Background Information**



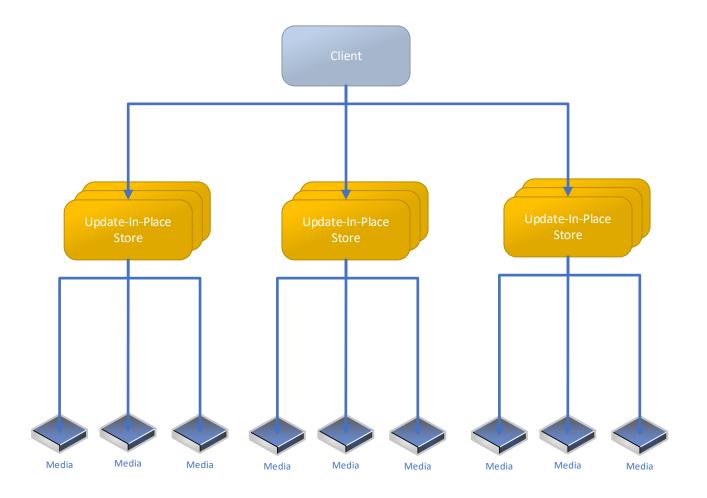
### Append-Only Store (Key/Value, Files, Etc.)



• Azure Storage ACM SOSP Paper (<u>link</u>), Presentation (<u>link</u>)



### Update-in-Place Store (Virtual Disks)



• Direct Drive - Azure's Next-generation Block Storage Architecture (link)



### **Indirection versus Redirection**

#### Indirection

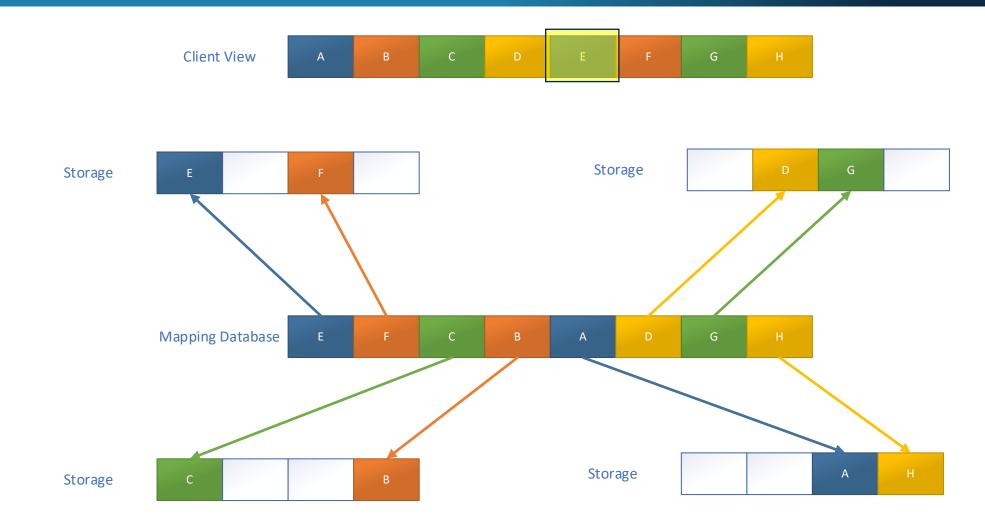
- Dynamic
- Requires additional storage and updates
- Requires a centralized "brain"
- Freedom to place data "anywhere"

### Redirection

- Static
- Just math
- Allows for shared understanding of data placement
- Forces data location

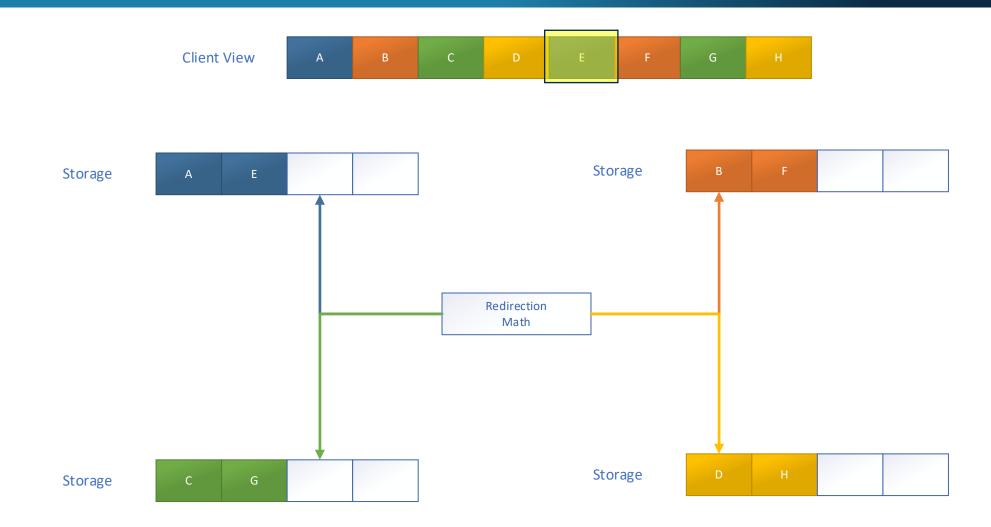


### Indirection





### Redirection



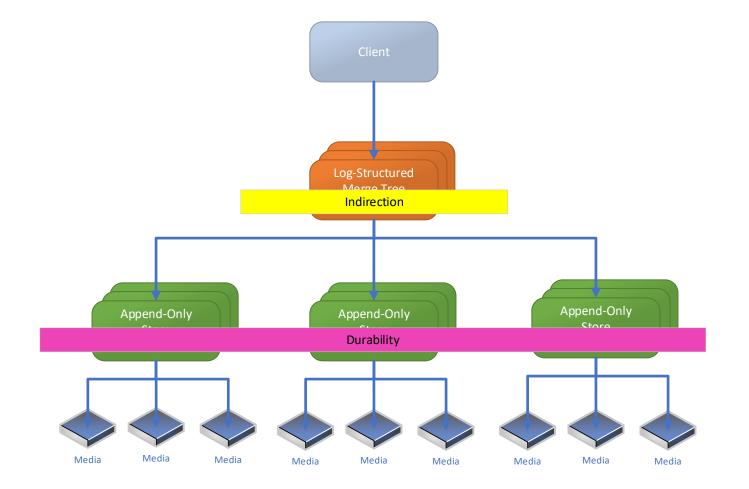




## Effects of Layering

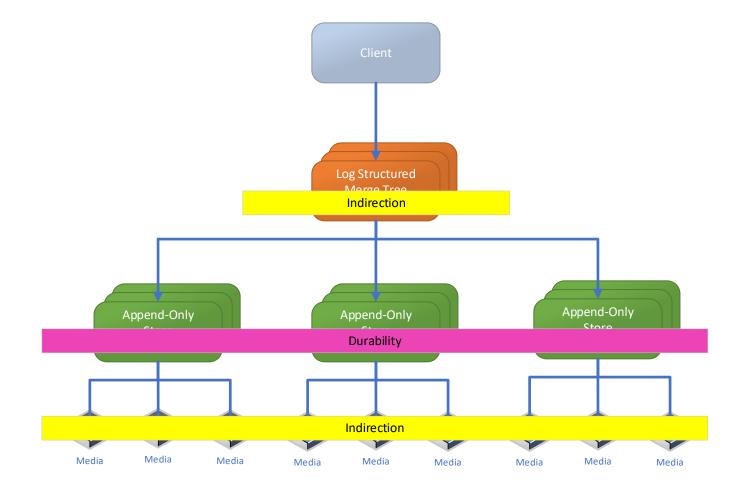


### **Indirection Above Durability**



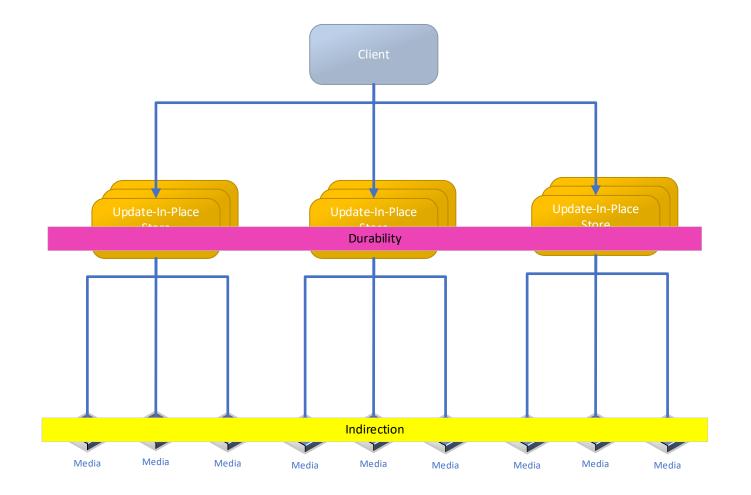


#### Rats...





### **Durability Above Indirection**





### **Indirection Location**

#### Indirection on top of durability

- Lends itself to append-only
  - Efficient snapshots
- Single "brain"
  - Reduces performance
  - Creates a single point of failure
- Easier costs reduction with things like erasure coding

#### Durability on top of indirection

- Lends itself to update-in-place
  - Inefficient snapshots
- Direct access to data provides highest performance
- Less flexibility of data placement

#### Multiple layers of indirection greatly increases WAF

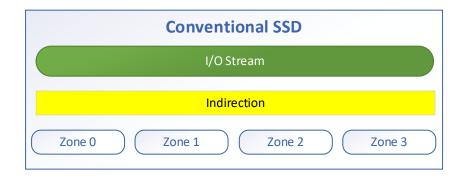


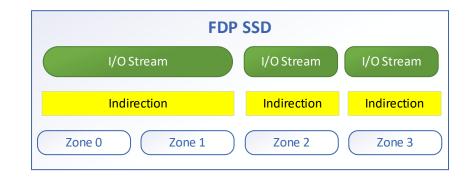


# New Media Types



### **Overview of Media**











### **Zoned Namespaces SSD**

#### Less RAM and overprovisioning

Saves cost

- Requires an indirection layer above it
  - Adds cost
- Scale is an important factor
- Allows for specialized indirection



## Shingled Magnetic Recording (SMR)

#### Very similar to ZNS

- Conventional area does not require indirection
  - No need for RAM
  - No need for overprovisioning
  - Allows for easier implementation of a traditional file system

#### Shingled area requires indirection\*

Major cost savings once you have this layer developed

\*Host managed assumed due to cost





# New Object Store



### **Specialized Object Store**

Provides both Update-in-Place and Append-Only Objects

### Built on ZNS and SMR

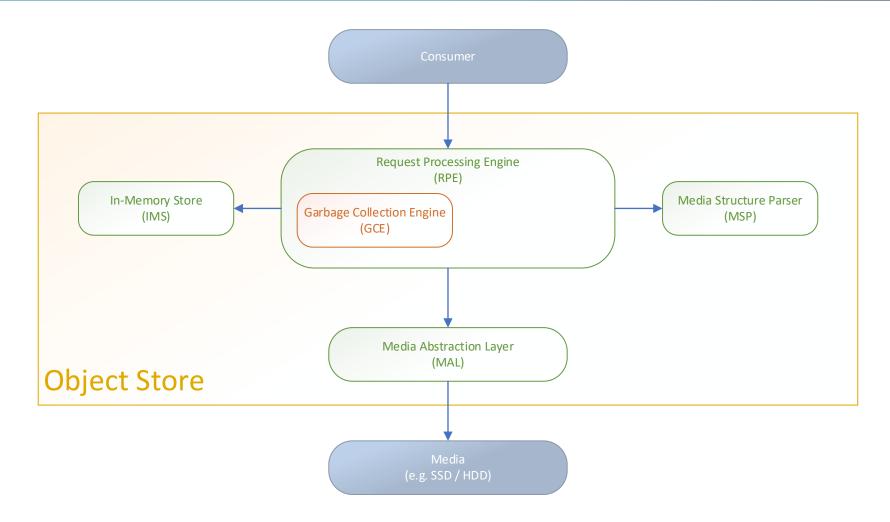
- Also works with conventional devices
- Abstracts media types

### Provides optional indirection on ZNS/SMR

- Reducing WAF requires a single layer of indirection
- Supports compression natively
- Minimal feature set
  - No namespace
  - Single consumer

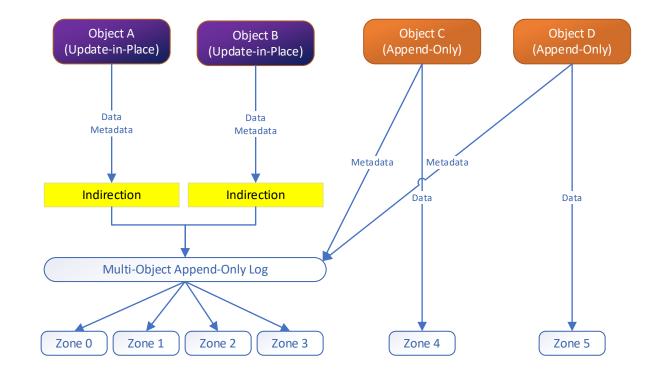


### **High-level Architecture**





### **Optional Indirection**







## **New Development Platform**



### Work Unit Based Development

- Inspired by DPU
- Similar to SPDK/DPDK
- Abstracts underlying platform
  - Allows for single code-base
  - Generates optimized code for each platform
    - Mutual exclusion can be implemented in different ways

#### Forces developer to break up logic into pipeline stages

- Tries to strike a balance between coarse vs fine-grain locks
- Single level of mutual exclusion means no deadlocks
- Designed for massive parallelism



### **Linear Coding**

```
WriteToObject()
{
    Lock(Object);
    CreateWriteRecord(Buffer);
    Lock(Device);
    Offset = Device->GetNextWriteOffset(Length);
    UpdateRecordWithOffset(Buffer, Offset);
    Device->Write(Buffer);
    Unlock(Device);
    UpdateObjectAfterWrite();
    Unlock(Object);
}
```



### **Linear Coding Improved**

```
WriteToObject()
{
    Lock(Object);
    CreateWriteRecord(Buffer);
    Lock(Device);
    Offset = Device->GetNextWriteOffset(Length);
    UpdateRecordWithOffset(Buffer, Offset);
    Unlock(Object);
    Device->Write(Buffer);
    Unlock(Device);
    Lock(Object);
    UpdateObjectAfterWrite();
    Unlock(Object);
```

}

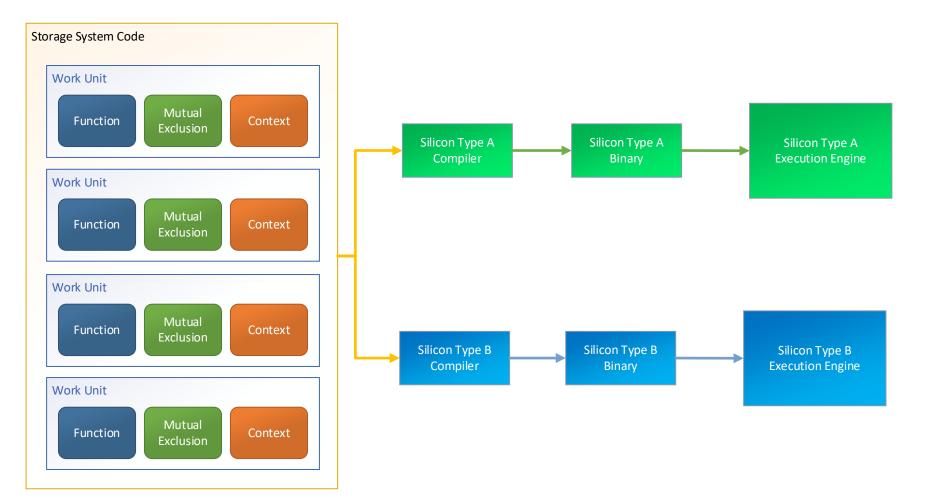


### Work Unit Based Coding

```
WriteToObject()
                                                    Write(Buffer, Offset)
  CreateWriteRecord(Buffer);
                                                        if(Offset == m_NextOffsetToWrite)
    ScheduleWork(UpdateRecordWithOffset, ...);
    ScheduleWork(Device->GetWriteOffset, ...);
                                                            DoWrites();
                                                        else
 UpdateRecordWithOffset(Buffer, Offset)
                                                            QueueWrite();
     Buffer->Offset = Offset;
     ScheduleWork(UpdateObjectAfterWrite);
     ScheduleWork(Device->Write, ...);
                                                      GetWriteOffset(Length)
          UpdateObjectAfterWrite()
                                                          *Offset = m_NextOffsetToUse;
                                                          m_NextOffsetToUse += Length;
              DoUpdates();
```



### **Run-to-Completion Platform**







## Conclusion

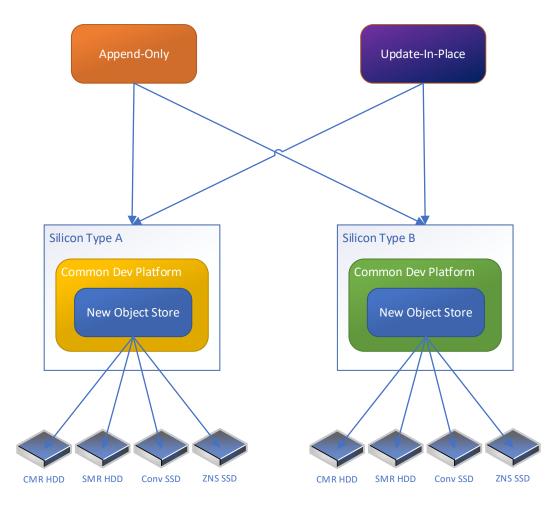




- Having separate systems provides robust product offerings
- Converging key areas reduces cost
- Abstraction of hardware reduces code duplication



### Vision for the Future







## Questions?





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