SNIA DEVELOPER CONFERENCE



September 16-18, 2024 Santa Clara, CA

High-Performance Block Volumes in Virtual Cloud Environment

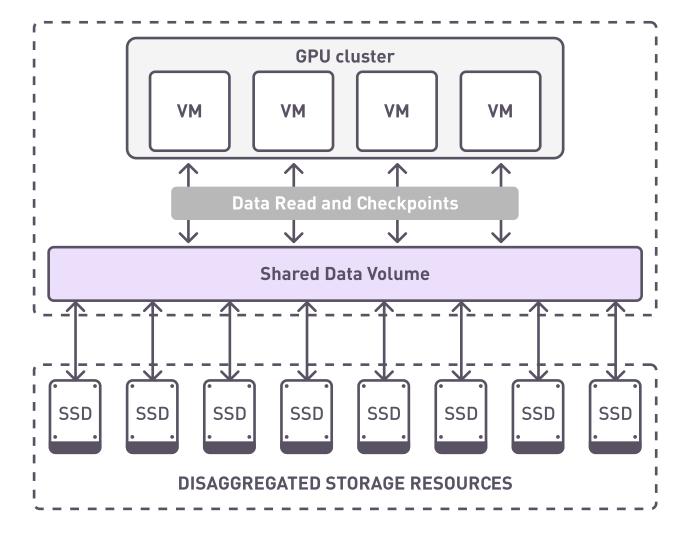
Solutions and Performance Analysis

Presented by Sergey Platonov and Heinrich von Keler

Al workloads in Cloud environments have a performance challenge



Problem: SDS performance with AI workloads in the Cloud

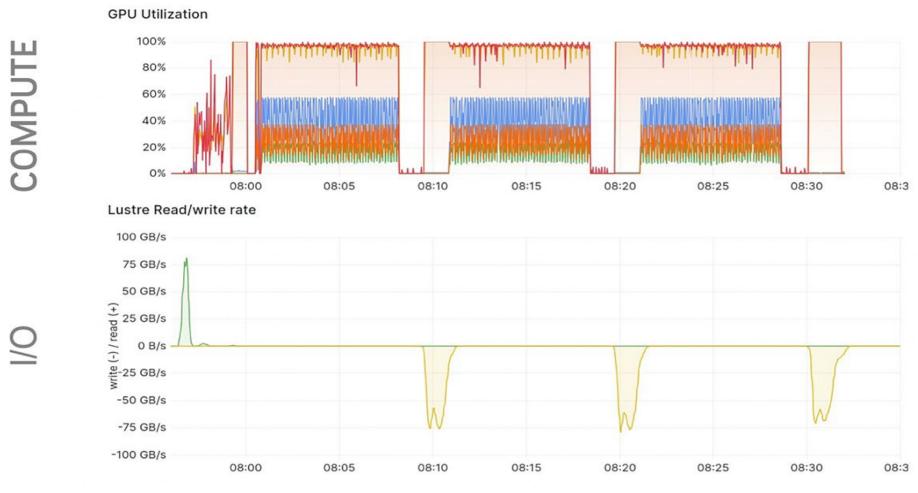


 Low Performance: SDSs often struggle to deliver the required performance for modern Cloudbased AI workloads.

Lack of Shared Volumes: SDS solutions rarely offer high performance shared volume support.



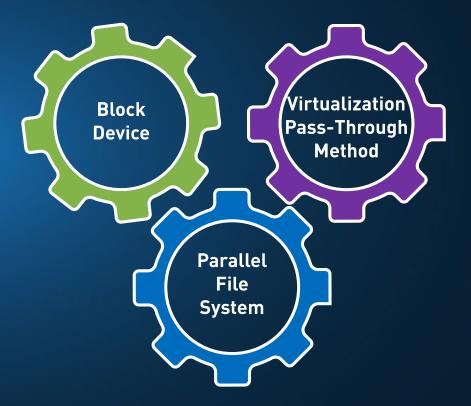
Slow Storage Means GPU Underutilization





Our Goals

- Develop performant block device.
- Block device must deliver highperformance when working under a virtualized parallel file system.
- Pass this solution through cloud environments without taking a performance hit with virtualization.
- Deliver random and sequential performance for AI workloads.







xiRAID Opus: optimal solution for block device in Cloud environments



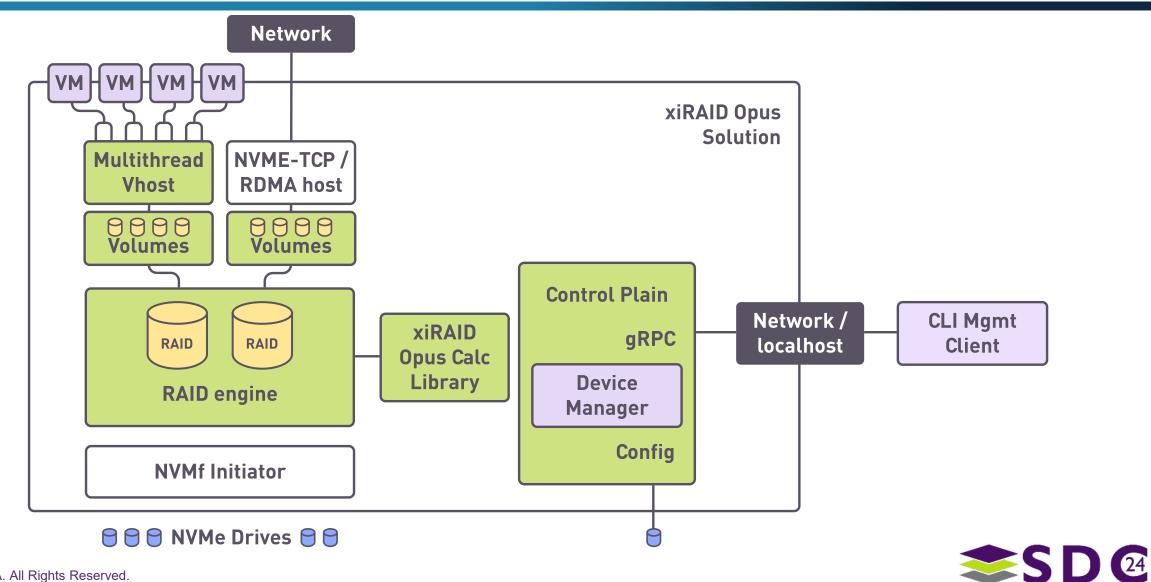
- Engine designed for high-performance storage data paths in virtualized environments.
- 1. Creation of RAID-protected volumes.
- 2. Provisioning of volumes to VMs.
- 3. Performance Enhancements:
 - Polling: Reduces latency by actively checking for I/O completions.
 - Zero-Copy: Eliminates unnecessary data copying, increasing throughput.

	Measured single drive performance	2x RAID5 theoretical performance	XiRAID 2x RAID5 performance	Efficiency
4K Random Read (M IOPS)	2,7	65	65	100%
4K Random Write (M IOPS)	0,7	8	8	100%
Sequential Read (GB/s)	14	336	310	92%
Sequential Write (GB/s)	6,75	149	144	97%

xiRAID demonstrates world-record performance with 24 Kioxia CM7 PCIe 5 NVMe SSD drives



xiRAID Opus Architecture



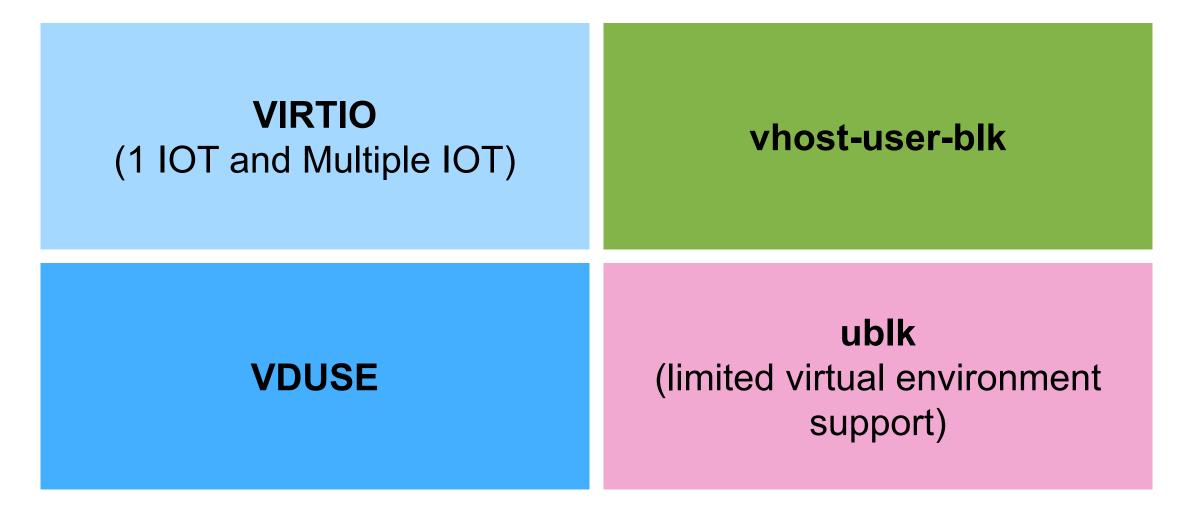


Most efficient method of block device delivery in virtualized environments.



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Methods for delivery of block device to VMs

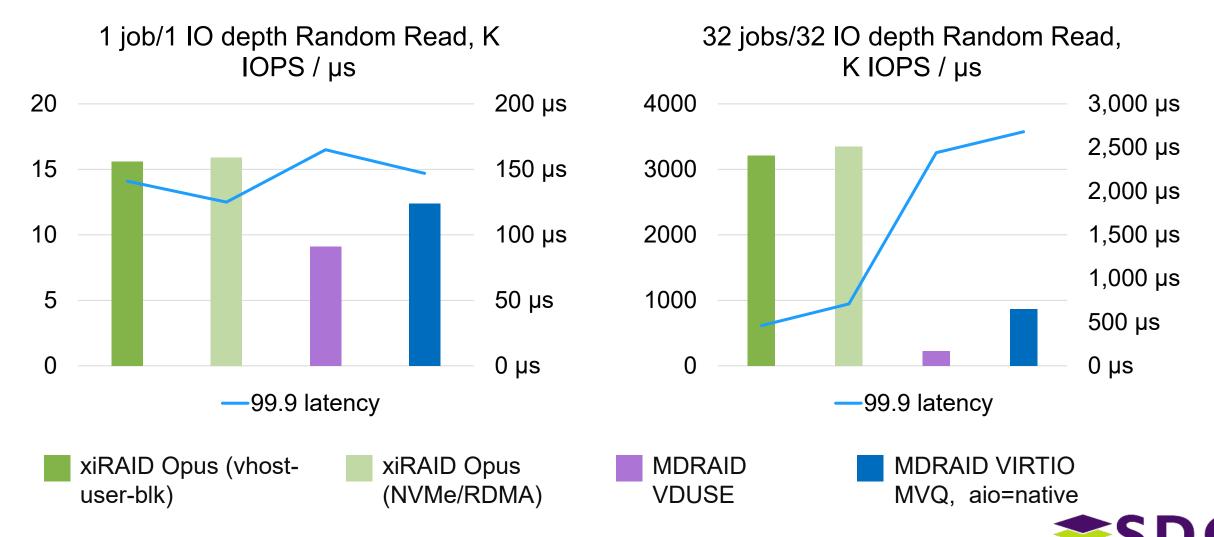




- xiRAID Opus RAID 5 (23+1 drives) vs MDRAID RAID 0 (24 drives)
- Single Virtual Machine: 32 VCPUs, 32GB of RAM
- Rocky Linux 9, kernel-lt(6.10)
- FIO v 3.36
- 4k random reads, AIO, direct_io
- full stripe writes, AIO, direct_io



Passing shared block volume to 1 VM - Random read

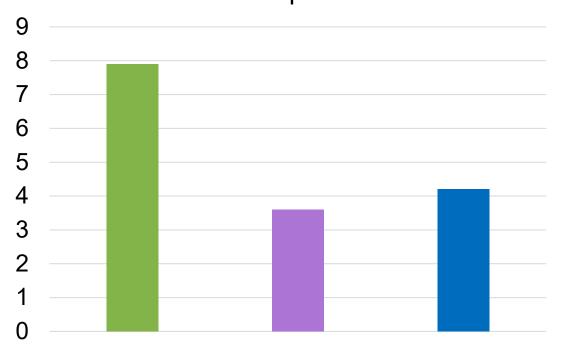


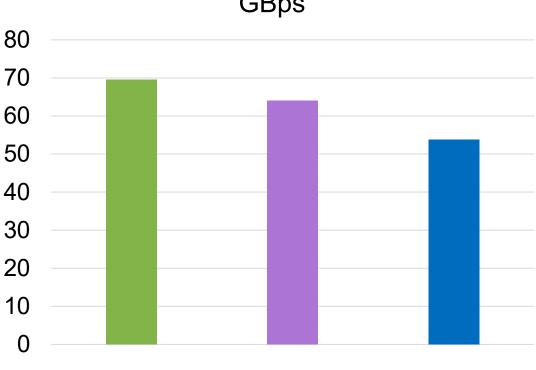
Passing shared block volume to 1 VM - Sequential write

MDRAID

VDUSE

1 job/1 IO depth Sequential write, GBps





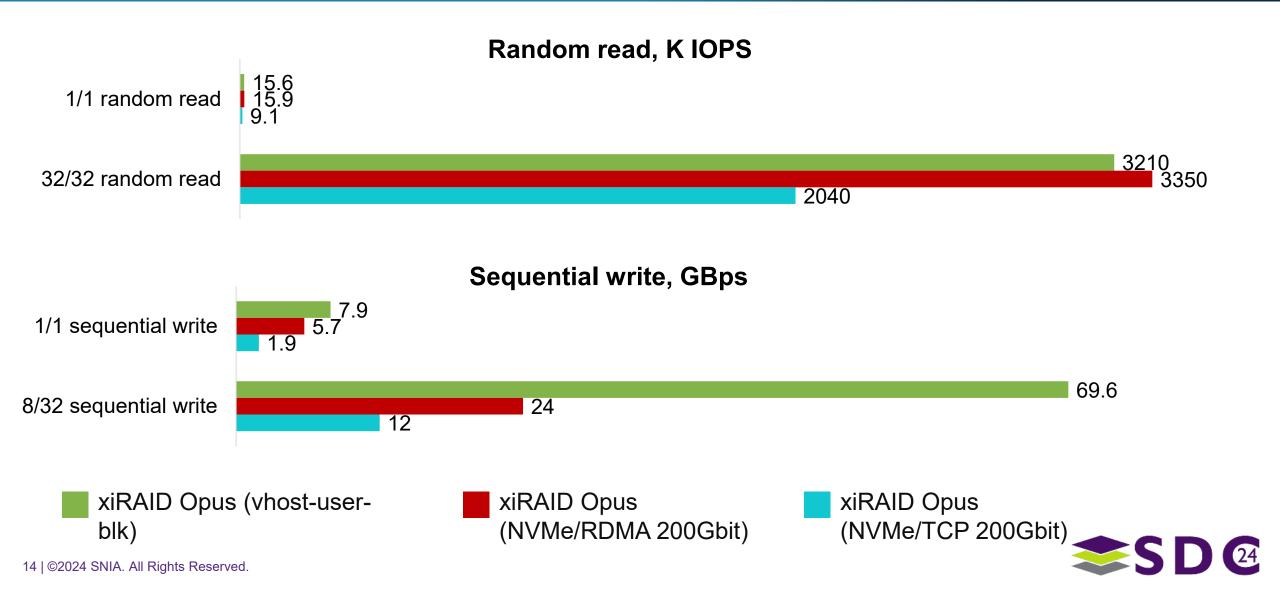
MDRAID VIRTIO MVQ, aio=native



xiRAID Opus (vhost-user-blk)

8 jobs/32 IO depth Sequential write, GBps

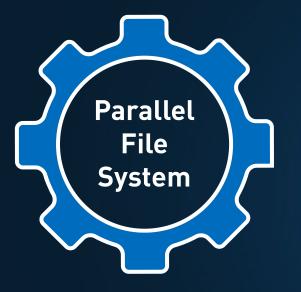
Passing shared block volume: Methods Comparison



xiRAID Opus using our implementation of vhost-usr-blk is the only option to deliver a high-performance block device to VMs.

- Developed multithreading interface for VHOST User BLK.
- Reused best practices from SPDK.
- Optimized utilization of specific CPU features from AMD and Intel.
- Optimized performance even for a single virtual machine and a single block device.





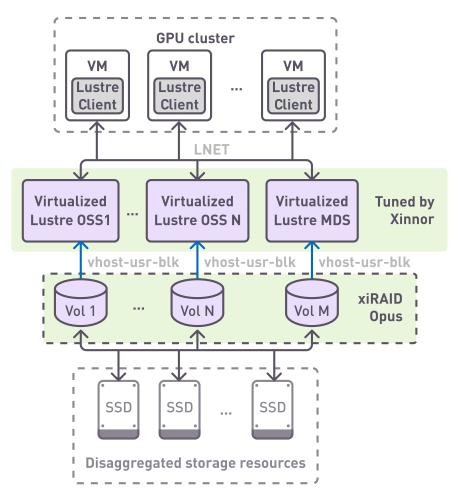
How xiRAID Opus accelerates parallel file systems in cloud environments



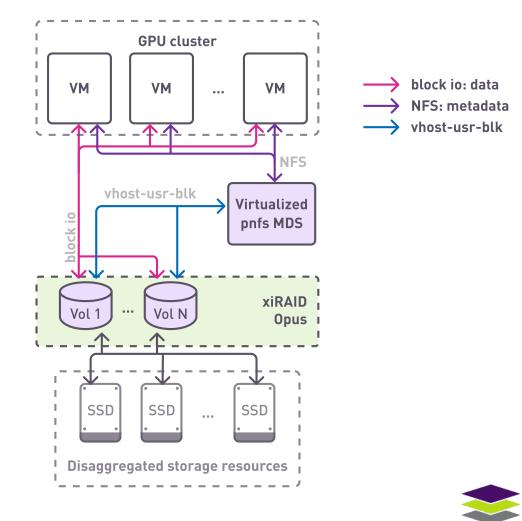
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Architectures that we are working on

1. Lustre solution



2. pNFS solution

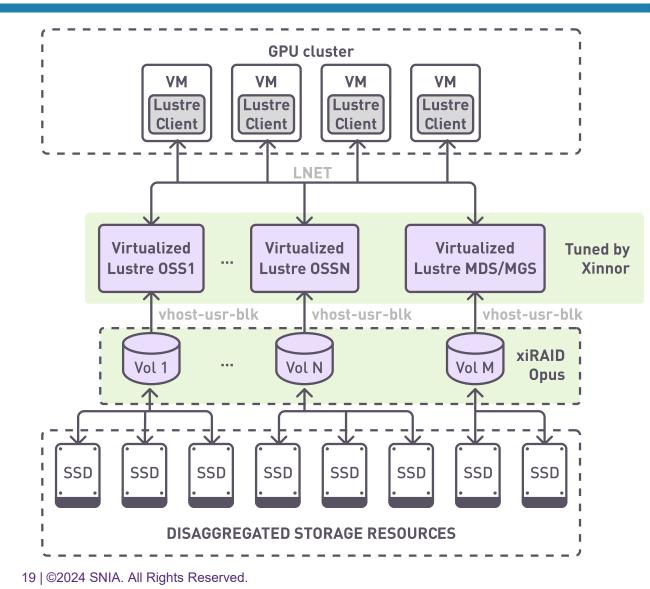


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Xinnor Lustre Solution for Cloud Environments

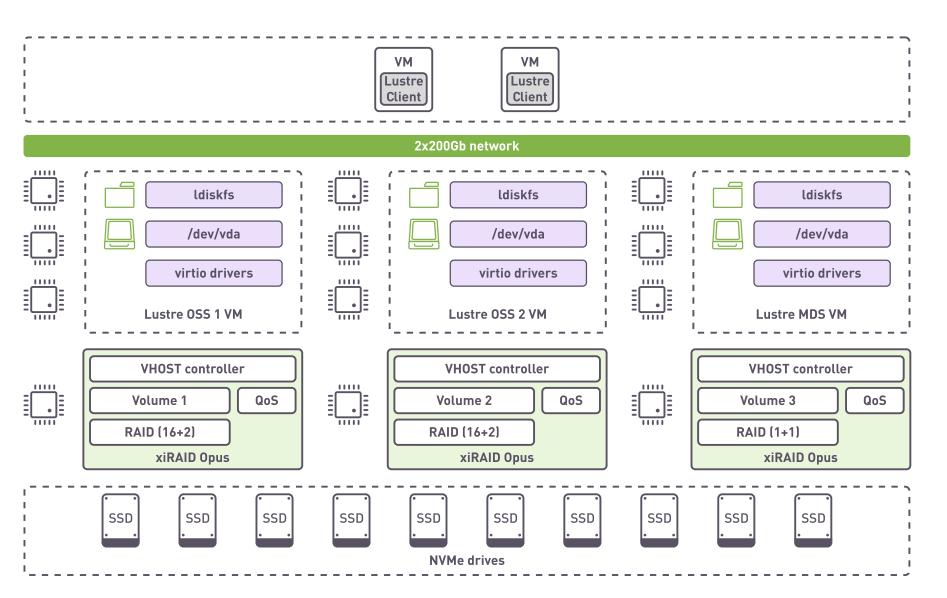


Lustre in Cloud Environments



- Lustre is a well–known FS that is primarily used for HPC workloads
- Provides good scalability and performance for data intensive workloads
- Provides HA over shared storage





Testing Environment Details

CPU: 64-Core Processor per node (AMD 7702P)

Memory: 256 GB RAM per Node

Networking: 1 x MT28908 Family [ConnectX-6] per node

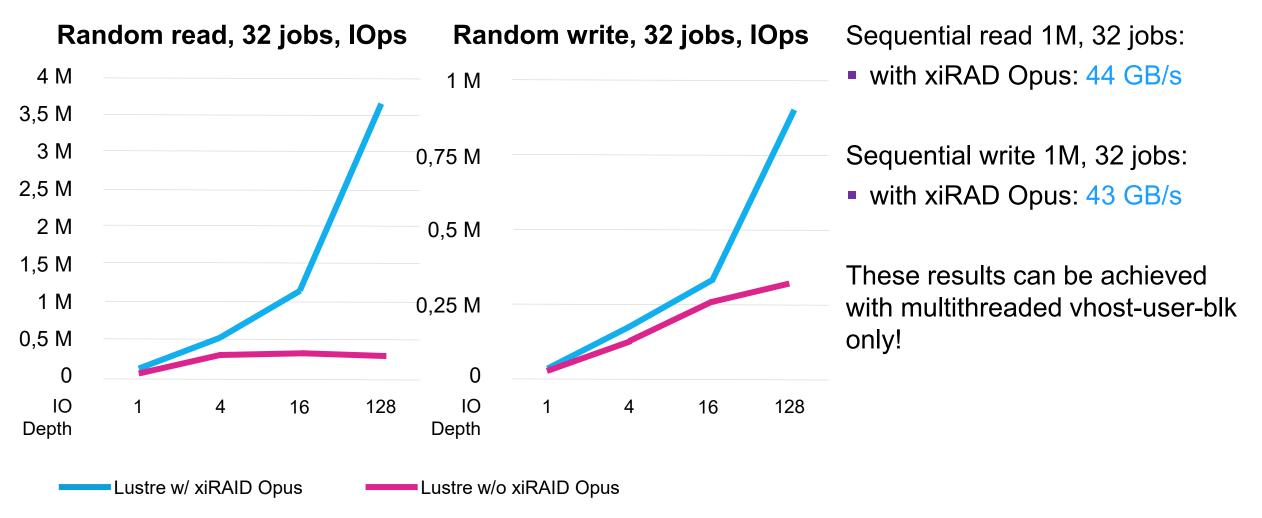
Drives: 24x KIOXIA CM6-R 3.84TB (Gen 4)

Aggregated drive performance per node:

- 9M IOps 4k RR
- 3M IOps 4k RW
- 70 GBps 128k SW, SR

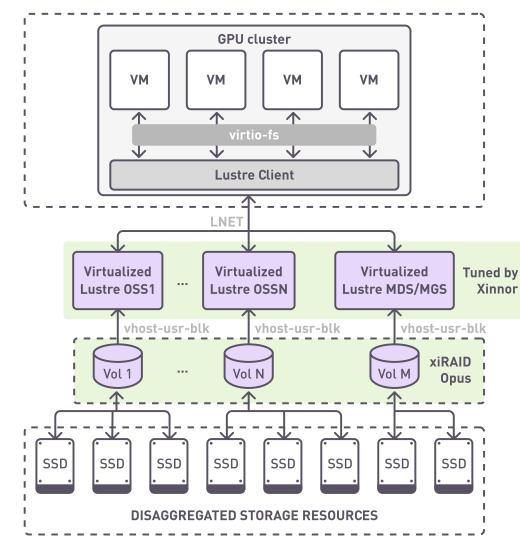


Lustre Solution Performance





Reducing complexity of Lustre administration

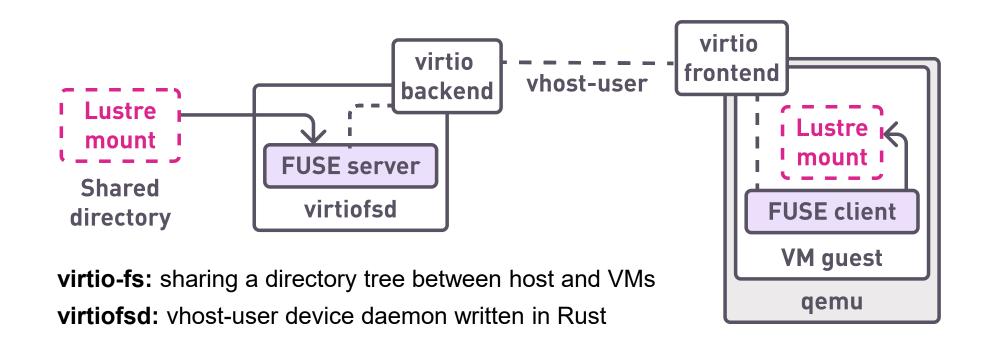


- Virtiofs allows to share mounted FS on the host with VMs
- No client software or specialized networking configuration needed



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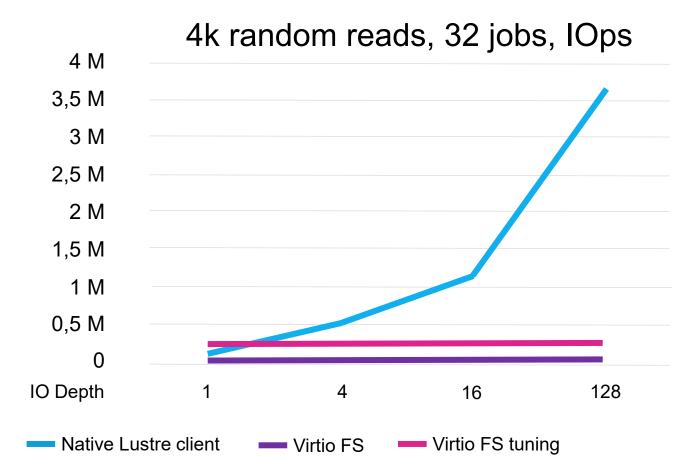
Virtio FS Upsides = simplicity



We can hide all the complexity of parallel file system setup from the client behind VIRTIOFS!



Xinnor Tuned Virtio FS = performance results



Sequential read 1M, 32 jobs:

- Native Lustre client: 44 GB/s
- Virtio FS: 9 GB/s
- Tuned Virtio FS: 44 GB/s

Sequential write 1M, 32 jobs:

- Native Lustre client: 43 GB/s
- Virtio FS: 7 GB/s
- Tuned Virtio FS: 44 GB/s



Outcomes = Xinnor Lustre Solution can perform

1. Performance:

- Even with only two virtualized OSS Lustre delivers strong results for sequential and random I/O operations (AIO).
- It is essential to have high-performance block devices passed through to the OSS and MDS virtual machines.
- xiRAID Opus solves this challenge.

2. Skill requirements:

- Requires a high level of expertise to configure the system and client VMs.
- VirtioFS can reduce complexity:
 - For use-cases that require only sequential workload patterns.

3. Xinnor can deliver Lustre solution for Cloud Environments.

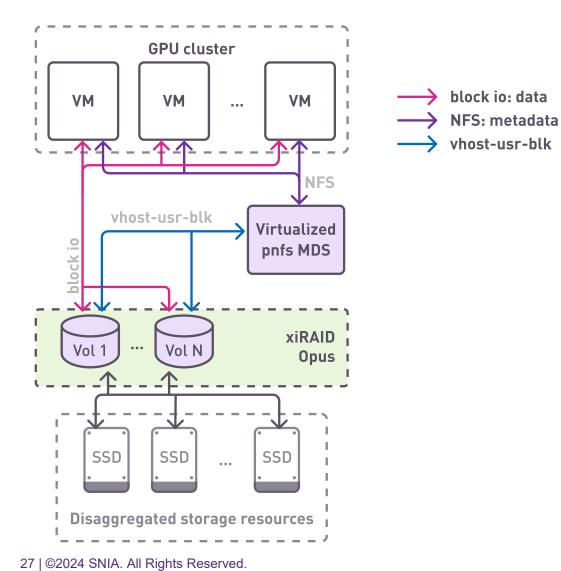
Reach out to Xinnor for a POC.



Our vision of the future: pNFS



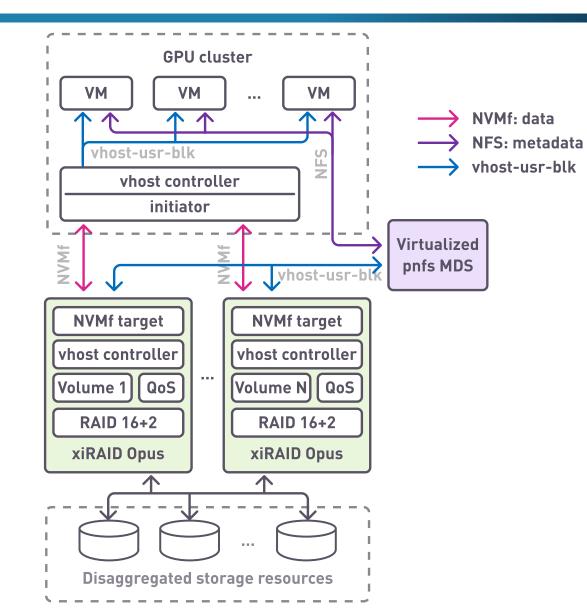
pNFS Block Layout



- Allows to build high-performance and scalable data volumes
- Clients reads and writes directly to volumes
- Data and metadata paths are separated
- No need for third-party client Software

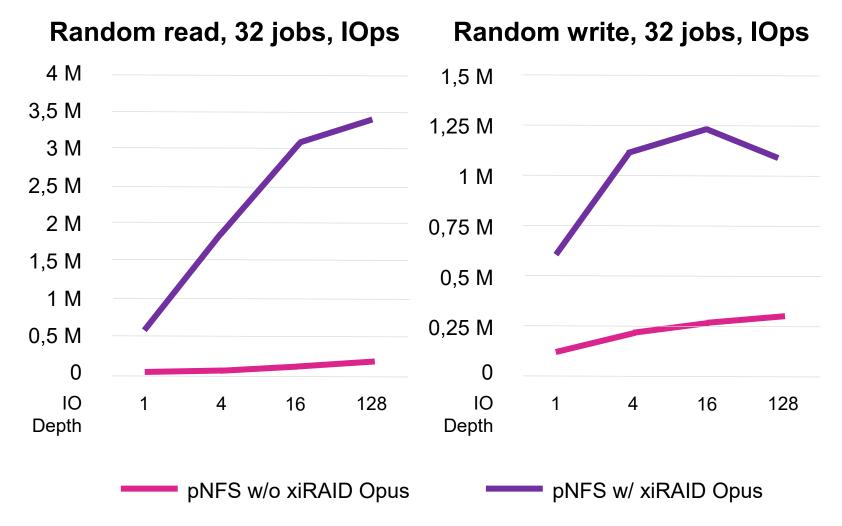


pNFS Architecture in Cloud Environments





pNFS with and without xiRAID Opus



Sequential read 1M, 32 jobs:

- without xiRAID Opus: 34,8 GB/s
- with xiRAID Opus : 47 GB/s

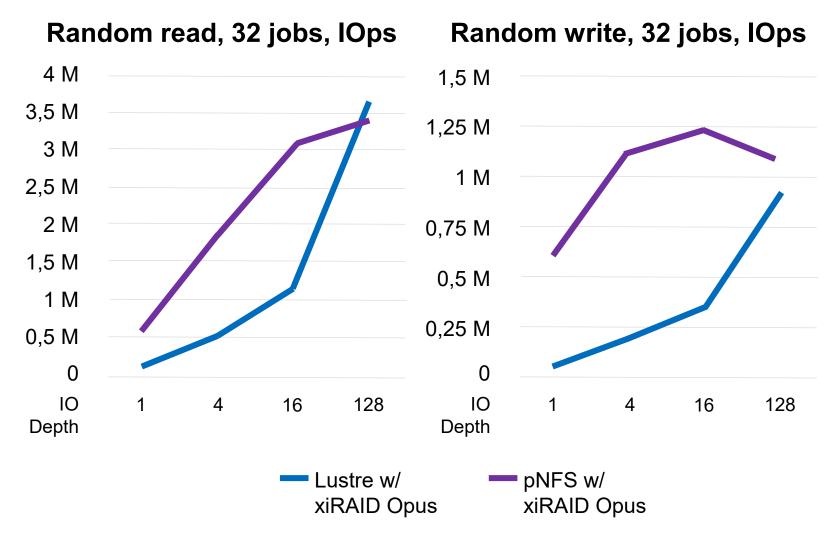
Sequential write 1M, 32 jobs:

without xiRAID Opus: 32,7 GB/s

with xiRAID Opus: 46 GB/s



pNFS vs Lustre - accelerated by Xinnor solutions



Sequential read 1M, 32 jobs:

- Lustre: 44 GB/s
- pNFS: 47 GB/s

Sequential write 1M, 32 jobs:

- Lustre: 43 GB/s
- pNFS: 46 GB/s



pNFS in Cloud Environments - Conclusions

Best Option with Proper Configuration

 Supports scalability to tens or even hundreds of gigabytes per second with minimal resource usage.

High Performance

 Sequential and random Small Block operations, with minimal latency due to direct interaction with the block device.

No 3rd party client software required

Drawbacks:

Open Source MDS is not production-ready and can only be used for POCs.

We are looking for partners to make pNFS production ready.



Conclusions



xiRAID Opus and Xinnor Lustre Solution are both ready to be deployed in Cloud environments as a high-performance solutions for Al workloads.

- Lustre Enables High-Intensity AI Workloads
 - New versions excel in Asynchronous Small Block I/O performance
- The VHOST-User-BLK interface, especially with Multi-IO Thread support, allows the direct passthrough of high-performance block volumes into virtual machines.







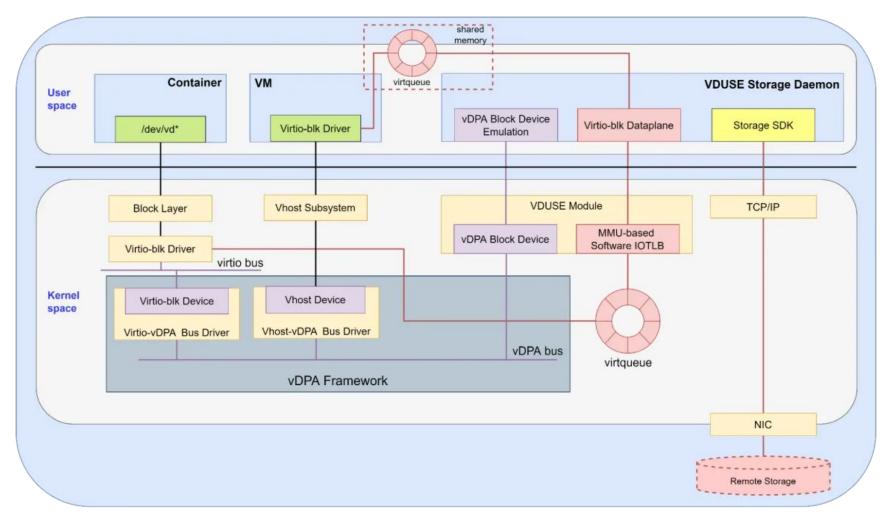
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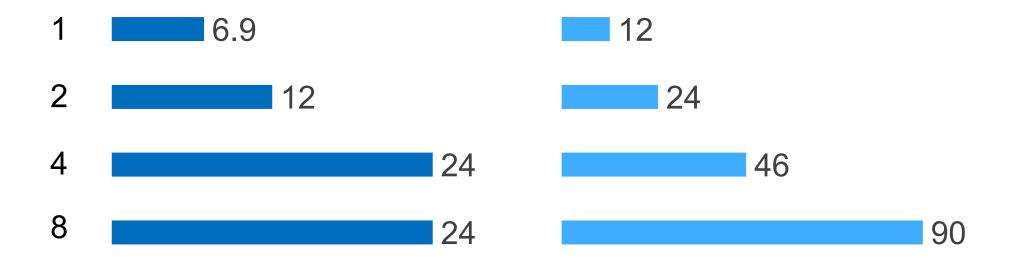
Backup





VirtioFS requires additional CPU cores for each VM

Thread pool size Sequential workloads, GBps Random workloads, K IOps





Outcomes of VirtioFS addition to Lustre

Simplified Configuration:

VirtioFS eliminates the need for client-side settings within virtual machines.

Performance Considerations:

- Tuning of VirtioFS on the host side is required.
- Dedicating extra CPU cores that are isolated from the ones used by the virtual machines is necessary.
- Increased cost of data access.

Random I/O Performance:

Poor small random I/O performance.

VirtIO FS is suitable for workloads that handle large volumes of data, where frequent Small IO operations are not required.

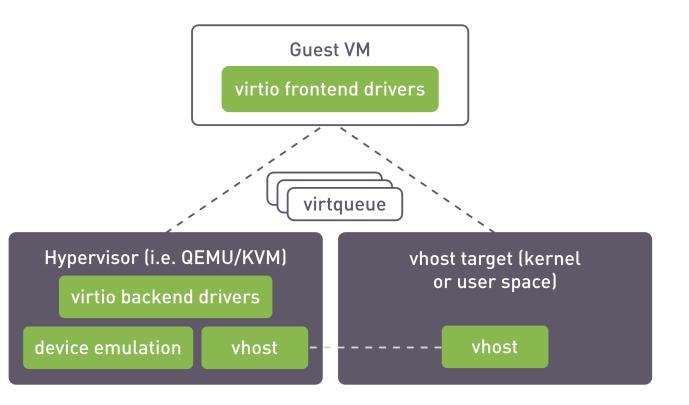
However, to achieve high performance, properly tuned Lustre OSS/OST and a high-performance backend are essential to efficiently handle the workload.



vhost-usr-blk

- Local block I/O interface
- Userpace
- Zero-copy (shared memory)
- Notifications and polling

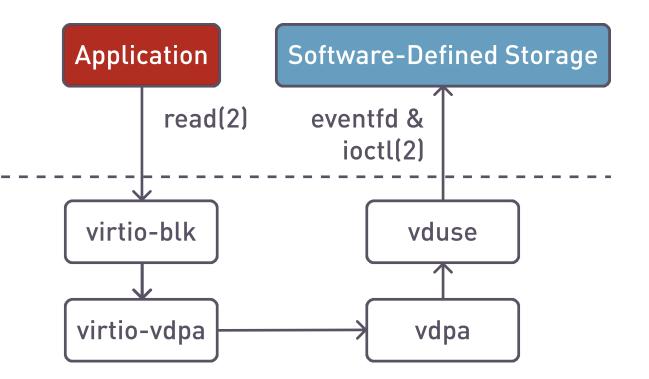
Linux, BSD, and macOS Implementations started in 2017





VDUSE: a software-defined datapath for virtio

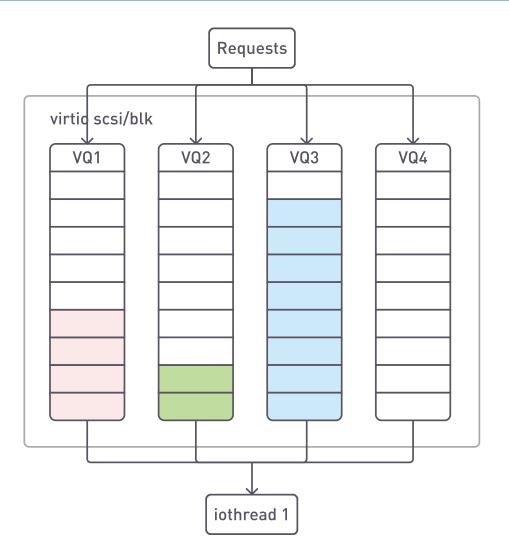
- vDPA device in userspace (<u>VDUSE</u>) is an emerging approach for providing software-defined storage and networking services to virtual machine (VM) and container workloads.
- The vDPA (virtio data path acceleration) kernel subsystem is the engine behind VDUSE.



VDUSE enables you to easily implement a software-emulated vDPA device in userspace to serve both VM and container workloads.

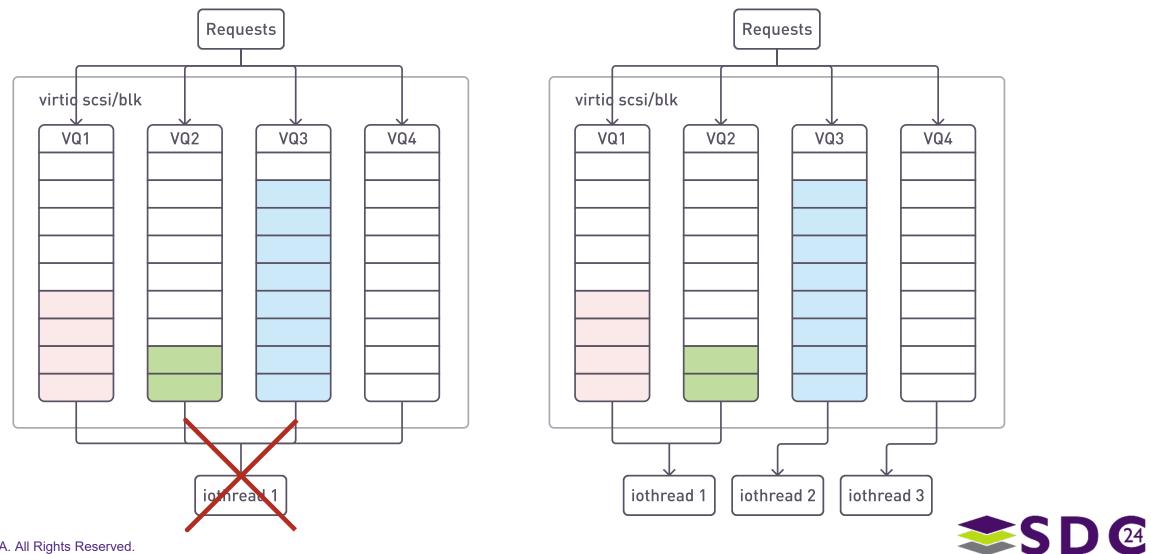


VirtIO - True multi-threading in QEMU 9



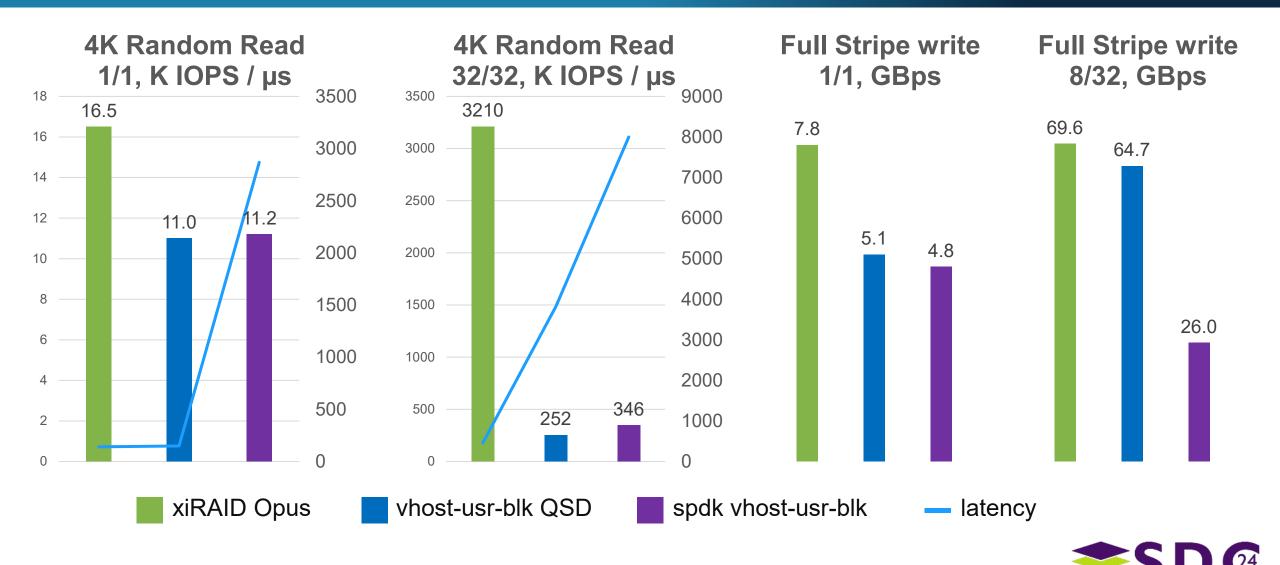


VirtIO - True multi-threading in QEMU 9



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Vhost-usr-blk Opus vs SPDK vs QSD implemetations



VIRTIO vs IO_URING PT

Workload	VIRTIO, MIOT, AIO=native		Io_uring PT	
Random read 1J/1IOD	12.4k IOps	147 µs 99.9 lat	12.7k IOps	133 µs 99.9 lat
Random read 32J/32IOD	870k IOps	2681 µs 99.9 lat	920k IOps	1647 µs 99.9 lat

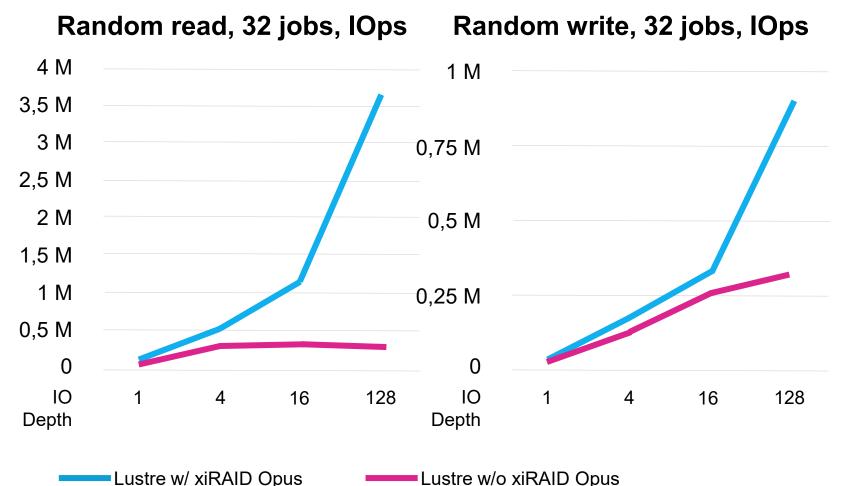


vDUSE, mdraid vs sigle drive vfio_pci

Workload	VDUSE 1 drive, vfio_pci		VDUSE, mdraid 0	
Random read 1J/1IOD	10.2k IOps	145 µs 99.9 lat	9.1k IOps	165 µs 99.9 lat
Random read 32J/32IOD	70k IOps	435 µs 99.9 lat	225k IOps	2440 µs 99.9 lat



Lustre Solution Performance



Sequential read 1M, 32 jobs:

without xiRAID Opus: 44 GB/s

with xiRAD Opus: 44 GB/s

Sequential write 1M, 32 jobs:

- without xiRAID Opus: 44 GB/s
- with xiRAD Opus: 43 GB/s

These results can be achieved with multithreaded vhost-user-blk only!

SD @

-name=randtest -bs=4k -ioengine=libaio -direct=1 -iodepth={1,32} numjobs={1,32} -- norandommap -- filename=/dev/vda -- group reporting -rw={randread,randwrite} -timebased=1 -runtime=600 -name=randtest iou -bs=4k -ioengine=io uring -fixedbufs=1 -hipri=1 registerfiles=1 -direct=1 -iodepth={1,32} -numjobs={1,32} norandommap -- filename=/dev/vda -- group reporting rw={randread,randwrite} -timebased=1 -runtime=600 -name=seqRW -bs={1024k, 1472k, 1536k} -ioengine=libaio -direct=1 iodepth={1,32} --numjobs={1,8} --offset increment=10% -filename=/dev/vda -group reporting -rw={read,write} -timebased=1 -

runtime=600



- -name=randtest -bs=4k –ioengine=libaio –direct=1 –iodepth={1, 8, 16, 128} numjobs={16} –norandommap -directory=/mount –group_reporting rw={randread,randwrite} -timebased=1 –runtime=600 –size=50G
- -name=seqRW -bs={1024k, 1472k, 1536k} –ioengine=libaio –direct=1 iodepth={1,32} –numjobs={1,8} -directory=/mount -group_reporting rw={read,write} -timebased=1 –runtime=600 –size=50G





taskset -ac 32-63 qemu-storage-daemon --blockdev driver=raw,nodename=md0,file.driver=host_device,file.filename=/dev/md0,cache.direct=o n --export type=vduse-blk,id=md0-export,nodename=md0,writable=on,name=vduse-0,num-queues=32,queue-size=256

taskset -ac 32-63 qemu-storage-daemon --blockdev driver=nvme,nodename=md0,driver=nvme,device=0000:01:00.0,namespace=2 --export type=vduse-blk,id=md0-export,nodename=md0,writable=on,name=vduse-0,num-queues=32,queue-size=256



taskset -ac 32-63 /usr/libexec/virtiofsd --socketpath=/tmp/vhostqemu.sock -o source=/virtiofs/ --cache=metadata --allowdirect-io --thread-pool-size=16 &



taskset -ca 0-31 /usr/libexec/qemu-kvm \

-enable-kvm -cpu host\

-m 32G -object memory-backend-file,id=mem,size=32G,mempath=/dev/shm,share=on -numa node,memdev=mem \

- -chardev socket,id=char0,path=/tmp/vhostqemu.sock2 -device vhostuser-fs-pci,queue-size=1024,chardev=char0,tag=myfs \ -smp 32 \
- -hda \${VMDISK_QCOW2} \
- -netdev user,id=net0,net=192.168.0.0/24,dhcpstart=192.168.0.9 \
- -device virtio-net-pci,netdev=net0 -vnc 0.0.0.0:2 --nographic



VM run parameters 2

taskset -a -c 0-31 /usr/local/bin/qemu-system-x86_64 \

-enable-kvm -cpu host \

-m 32G -object memory-backend-file,id=mem,size=32G,mem-path=/dev/shm,share=on -numa node,memdev=mem \

-smp 32 \

-hda \${VMDISK_QCOW2} -drive if=none,id=drive0,cache=none,aio=native,format=raw,file=/dev/md127 \

-device '{"driver":"virtio-blk-pci","drive":"drive0","iothread-vqmapping":[{"iothread":"my0","vqs":[0,1]},{"iothread":"my1","vqs":[2,3]},{"iothread":"my2","vqs":[4,5]},{"iothread":"my3 ","vqs":[6,7]},{"iothread":"my4","vqs":[8,9]},{"iothread":"my5","vqs":[10,11]},{"iothread":"my6","vqs":[12,13]},{"iothread d":"my7","vqs":[14,15]},{"iothread":"my8","vqs":[16,17]},{"iothread":"my9","vqs":[18,19]},{"iothread":"my10","vqs":[20,21]},{"iothread":"my11","vqs":[22,23]},{"iothread":"my12","vqs":[24,25]},{"iothread":"my13","vqs":[26,27]},{"iothread":"my14","vqs":[28,29]},{"iothread":"my15","vqs":[30,31]}]}' \

-netdev tap,id=net0,ifname=tap0,script=no,downscript=no -object iothread,id=my0 -object iothread,id=my1 -object iothread,id=my2 -object iothread,id=my3 -object iothread,id=my4 -object iothread,id=my5 -object iothread,id=my6 -object iothread,id=my7 -object iothread,id=my8 -object iothread,id=my9 -object iothread,id=my10 -object iothread,id=my11 -object iothread,id=my12 -object iothread,id=my13 -object iothread,id=my14 -object iothread,id=my15 \-device virtio-net-pci,netdev=net0 -vnc 0.0.0.0:1



taskset -ac 0-31 /usr/local/bin/qemu-system-x86_64 \

-enable-kvm -cpu host -m 32G -smp 32 \

-m 32G -object memory-backend-file,id=mem,size=32G,mempath=/dev/hugepages,share=on -numa node,memdev=mem \ -hda \${VMDISK_QCOW2} \

-vnc 0.0.0.0:2 \

-device vhost-vdpa-device-pci,vhostdev=/dev/vhost-vdpa-0,queuesize=256



taskset -a -c 0-31 /usr/libexec/qemu-kvm \

-enable-kvm -cpu host\

-m 32G -object memory-backend-file,id=mem,size=32G,mempath=/dev/hugepages,share=on -numa node,memdev=mem \ -smp 32 \

-hda \${VMDISK_QCOW2} -chardev socket,id=char1,path=/opt/xiraid/bin/xnr_conf/sock/r1 - device vhost-user-blk-pci,id=blk0,chardev=char1,num-queues=32,queue-size=256 \ -netdev socket,id=net0,connect=127.0.0.1:1234 \

- -device virtio-net-pci,netdev=net0 \
- -netdev user,id=net1,net=192.168.0.0/24,dhcpstart=192.168.0.9 \
- -device virtio-net-pci,netdev=net1 \
- -vnc 0.0.0.0:4 --nographic

